

CICS TS Performance Tuning Tutorial

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- YMMV
- Remember the Political Factor
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 - Introduction
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Introduction

- CICS uses three techniques to handle VSAM files within CICS TS:
 - Non-Shared Resources (NSR)
 - Local Shared Resources (LSR)
 - Record Level Sharing (RLS)
- In recent years, new VSAM features announced for CICS have been LSR oriented
- The major difference between the three techniques lies in the "ownership" of the resources
 - NSR \rightarrow resources are used exclusively by the file
 - LSR → resources are shared between participating files
 - RLS → resources are "owned" by a separate address space (SMSVSAM)



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Introduction

I/O generates CPU usage

- Application request to
- CICS FC programs to
- VSAM to
- SVC Handler to
- IOS
- Start the I/O (SSCH) and eventually back to
- CICS to have task wait
- Process I/O Interrupt
- Create SRB
- Dispatch the SRB to Post Completion
- To the CICS Dispatcher that dispatches the task when its turn occurs
- To improve response time and reduce CPU overhead, you need to eliminate I/O
 - Find the data/index in a buffer called a Look-Aside Hit
 - CPU requirements for a Look-Aside Hit is much lower
 - Return to CICS and application come from VSAM





The Very Big I/O Picture





Non-Shared Resources

NSR





Introduction to NSR

- NSR advantages include:
 - Resources are reserved so one file can be specifically tuned
 - Allows for chained read operations that can give better sequential performance
 - BROWSE
 - CA Splits
 - Mass inserts
- Does not support Transaction Isolation
- Does not support VSAM Threadsafe
- NSR = BATCH Processing





NSR File Definition

- A file is defined as NSR by specifying LSRPOOLNUM (NONE)
- String number defines the number of concurrent file accesses allowed
- One BUFND and one BUFNI (if applicable) is required per string
- Minimum buffer allocations:
 - BUFND is string number plus one
 - Extra buffer is only used for split processing
 - BUFNI is string number



NSR File Definition



- String definition for an NSR file can be a challenging task
 - Many NSR files are over allocated in strings when considering the I/O activity against the file
 - The major reason is that NSR allows duplicate CIs to exist between strings
 - NSR allows STARTBR/READNEXT/READ for UPDATE sequence without an intervening ENDBR
 - This results in two strings being allocated to the task
 - The requested CI appears 2X in VS
 - As a result, many files would appear to be deadlocked due to lack of strings
 - This type of request will not work in LSR
 - Remember that a string needs a BUFND/BUFNI
 - Eliminate strings in favor of more buffers





NSR File Definition

- Additional buffers can be allocated
 - Extra BUFND will be used in sequential operations
 - All available buffers will be allocated to the 1st sequential request
 - Extra BUFNI will be used to store Index Set (IS) indices (high level indices)
 - Sequence Set Indices (SSI) are never read into the extra BUFNIs
 - SSI CIs are read into the string index buffer
 - No look aside to other string buffers are done



NSR Buffer Definition



• Example # 1:

• STRNO = 2 BUFND = 3 BUFNI = 2





NSR Buffer Definition



• Example # 2:

• STRNO = 2 BUFND = 4 BUFNI = 3









• Why would you want extra BUFNDs?

- In the case of a BROWSE request, to read ahead a number of CIs to improve performance of the task
- In the case of Mass Insert, to write behind a series of CIs to improve task performance
- In the case of CA Splits, to be able to move more than one CI at a time to the new CA
- Overall, extra data buffers can speed up the process and reduce I/O requests to the file







What is the hidden agenda?

Browse

- The number of BUFNDs defined should contain the approximately the same number of records read (READNEXT) by the program
 - For example, if a CI can contain 5 records and the average # of READNEXT operations issued is 20, then a BUFND specifying 4 additional buffers (5 records/CI*4 read ahead buffers) would be fine
 - However, what programmer knows on the average how many READNEXT operations are issued to a file?
 - Also, only the 1st BROWSE request would benefit
 - What happens if the BROWSE is ended (ENDBR) before the 20 READNEXT operations are done?
- Adding additional buffers for sequential BROWSE processing will increase the task response time plus elongated I/O operation will result
- In addition, having the data in storage is good for this task but may affect the response of other tasks in the system





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- Mass Inserts
 - The number of buffers should be around the same number of writes (WRITE) issued to the file at one time
 - Same logic as the BROWSE
 - However, if the number of writes ends before all the buffers are full, then there is no I/O penalty as in the case of a BROWSE
- CA Splits
 - The number of buffers should be large enough to copy ½ of a CA at time
 - However, if the file does Mass Inserts or BROWSE operations, there is no way to segregate the buffers for one particular use







- What is the best approach for files that are heavily or mainly browsed?
 - If too many buffers are read, performance of other tasks may be affected
 - The key is to try and get a CISZ that generally accommodates the # of READNEXT commands issued
 If too many, try to get a large multiple
 - This approach can be used for LSR pool files too







- Why would you want extra BUFNIs?
 - To improve the look-aside and reduce physical I/O
 - Two types of index look asides occur for an NSR file
 - The 1st look aside is for the Index Set records that are in extra BUFNI buffers
 - The 2nd look aside is within the string buffers to see if the Sequence Set Index and/or the data CI are present
 - No look aside possible to other string buffers







- Additional index buffers allows VSAM to load the Index Set records
 - User should allocate sufficient BUFNIs as there are Index Set CIs in the file
 - Consideration should be given to adding additional buffers if the file reflects CA splits
 - Data CA splits can cause index CA splits creating new index set records





- Determining the number of BUFNIs required entails computing how many Sequence Set Index (SSI) records exist in the file
 - There is one Sequence Set record per data CA
 - This is a one to one relationship





• Compute:

- 1)# CAs = (Data HURBA / (# CI/CA*Data CISZ) this represent the # of Sequence Set Index records in the file
- 2)From LISTCAT get the total number of Index records in the file and determine the number of Index Set records in the file: (Total Number of Index Records – # of CAs)
- 3) Determine the # of BUFNIs = (Total # Of Index Set records + # of strings + CA split adjustment)
- 4)CA Split adjustment is any figure from zero to "n", where "n" is the # of additional Index set records created as a result of CA splits





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• Example using previous LISTCAT information

2

- Data CISZ 18K (18,432)
- CI/CA 45
- Bytes/CA 829,440 (18432*45)
- CA splits Yes
- # of IX records 4
- HURBA 2,488,320
- # of IX levels
- (2488320/829440)=3 CAs or Sequence Set Records
- (4-3)=1 Index Set Record
- If STRNO=5, then (5+1+2)=8 BUFNI request for the file. The +2 is a buffer for future CA splits at the index level. The CA adjustment is optional and the value can vary



NSR Buffer Definition

- Example # 1 VSAM 2 Index Levels:
 - STRNO = 2 BUFND = 3 BUFNI = 2
 - Requires three I/Os (2 index and 1 data)
 - No opportunity for look aside









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NSR Buffer Definition

- Example # 2 VSAM 2 Index Levels:
 - STRNO = 2 BUFND = 4 BUFNI = 3
 - After 1st read, each request would require a maximum of two reads or a 33% I/O operations savings





NSR Recommendations

- NSR files should be reviewed to see why they are not in LSR for better performance
 - For example, Share Options 4 file
 - Command Level Browse restrictions
- If the file is to be in NSR
 - Ensure valid CISZ for files that are browsed
 - Ensure sufficient BUFNIs allocated to hold the entire Index Set indices in buffers
 - Ensure that excess strings are eliminated and the storage used to allocate correct file buffering
 - Do not over allocate BUFND unless the file is prone only to CA splits
- If NSR must be used and files takes CA splits, consider activating the CO TCB (SUBTSKS=1 in SIT)
- NSR and Transaction Isolation are incompatible
- NSR is not supported under VSAM Threadsafe





Local Shared Resources

LSR





Robin Hood Theory

- Tuning LSR files is simply applying the Robin Hood theory in reverse
- In Sherwood Forrest Robin stole from the rich to give to the poor
- In LSR you steal from the poor to give to the rich!!!!!
 - **Poor = Low to Medium Activity Files**
 - Rich = Most Active Files
- In other words the major contribution that low activity files provide to LSR are their resources so that higher activity files can use them (Cruel Reality)





Introduction to LSR

• LSR advantages include:

- More efficient VS use because resources are shared
- Better look-aside because index buffers can maintain the Sequence Set Index records
- Tends to be more self-tuning because buffers are allocated on an LRU basis keeping information of the more active files in the buffers at the expense of less active files
- Only one copy of a CI allowed (better read integrity)
- Can allocate up to 255 pools to segregate files
- Supports Transaction Isolation (TI)
- Supports VSAM Threadsafe (Local VSAM)





LSR Tuning Areas

- Pool definition dynamic or static
- Buffer hit ratios
- Buffer monopolization
- Number of LSR pools
- Overlooked tuning opportunities
 - Buffer fragmentation
 - Buffer vs. CISZ
 - Page allocation
 - Maximum key size
 - Number of strings
- LSR candidates



Dynamic vs. Static Pool Definition



Dynamic Definition

- Advantages
 - Easy to implement
 - Little SysProg intervention
- Disadvantages
 - Combined buffer pool for data and index
 - Resource allocation based on a percent
 - Slow initialization
 - Cannot tune specific buffer sizes



Advantages

- Separate buffer pool for data and index can be defined
- Resource allocation can be optimized by activity
- Faster initialization

Disadvantages

•

- Requires SysProg intervention
- Can be prone to errors
- Requires planning



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Pool Definition

- Recommendation
 - Define LSR Pools Explicitly
 - Determine Individual File Requirements
 - Data and Index (If Applicable) CISZ required
 - Maximum Length Key
 - Strings
 - Get "Big Picture" of Requirements
 - CICS Performance Tool/Monitor
 - CICS Statistics (EOD)
 - Dynamic Definition One Time





- LSR pool effectiveness is based on look-aside hit ratios
 - Generally accepted hit ratios ()
 - Data 80%+
 - Index 95%+
 - Combined 93%+
- Buffer tuning should concentrate on improving the index hit ratio first
 - Generally, index I/O is higher than the data
 - Virtual and real storage investment to improve index hit ratio is less due to smaller CISZ associated with the index component





- Important note:
 - LSR buffer attainments can be misleading
 - If the 4 KB buffer reflects a hit ratio of 85%, this does not mean that every file is getting an 85% lookaside hit ratio
 - The 85% is an average of all the files using this buffer size
 - Some get a higher attainment
 - Others get a lower attainment



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- Data buffer tuning is highly dependent on access patterns
 - Good look-aside hit ratios usually requires a substantial storage investment (80%+)
 - The major cause Is that the data component is usually very large (vs. index component)
 - Good hit ratios usually result in files with:
 - Sequential activity
 - Read for Update/Rewrite/Delete
 - Concentrated read activity



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- Data buffer tuning is highly dependent on access patterns
 - Bad hit ratios usually result in files with:
 - Disperse read activity (very large files)
 - Share Options 4
- Recommendation
 - Buffer tuning is usually a "trial and error" process in determining the number of buffers to add to each buffer size
 - Reiterative process
 - You add buffers
 - You measure
 - If objective met, temporary end, else go back to add buffers
 - Temporary end because things change and require periodic observation
 - Tune buffer pools and CI sizes individually
 - Set Realistic Objectives, for Example:
 - Data 80%
 - Index 95%
 - **Combined 93%**
 - Define a minimum of three 32K catch-all buffers or both the data and index component







Buffer Monopolization

Buffer monopolization

- "Monopolization" refers to the buffer size within the LSR pool
- Theory behind LSR is to share resources when needed
 - So what can be bad if the principal files (most active) control a high percentage of the buffers?
 - Even at the expense of low activity files
- How do you determine if a file is monopolizing a particular buffer size?
 - I/O activity
 - Buffer hit ratio
 - Number of buffers held (by CISZ)





Buffer Monopolization

Buffer pool monopolization

- Need a CICS tuning/monitor to determine the number of buffers being held by a file
- Important if principal files are not providing a good response time
- Remember the reverse "Robin Hood Theory"
 - "Rob from the poor to give to the rich"
 - Where the "rich" are your most important active files
- Point of Diminishing Return
 - Keep adding buffers until the higher activity files do not require more





Number of defined LSR pools

- Two schools of thought
 - School 1 Use as many pools as possible so that files can be segregated to reduce buffer contention and/or interference
 - School 2 Use as few as possible pools so that resources can be used more efficiently
- Considerations
 - Are the pools allocated with a "Fudge Factor"?
 - Which files are more important so that resources should be allocated to them?





- There are 255 (MVS/CICS TS V4.R2) (originally 8) or 15 (VSE) LSR Pools available
 - The use of multiple buffers made sense with the original release of LSR support in CICS because the buffer search algorithm was sequential
 - Larger pools increased CPU time to search
 - Breakeven point was around 200 buffers
 - Another major benefit for LSR was virtual storage savings vs. NSR
 - Search algorithm changed to hashing technique
- The Least Recently Used (LRU) algorithm used by VSAM is self-tuning
- Access to the pools was single threaded on the QR TCB
 - Having multiple pools did not mean that there could be any overlapped access to different pools









- There may be some cases where you may want to define extra pools
- Data Tables
 - Output operations go against the VSAM file
 - LSR pool used for look-aside for records before going to disk
 - ROT for Data Tables = 90%+ Read Operations
 - Low activity to the pool will reduce look-aside capacity
 - Define a separate pool for all data tables giving more buffers to the index component
- Favor one or more File
 - Important file(s) that you want to give special attention and do not want any contention for buffers or strings
- A pool that needs more than 255 strings





• LSR VSAM Threadsafe files

- Allows for the use of Open TCBs to handle VSAM requests
 - Prior to CICS TS V4.R2 you had 8 LSR pools
 - CICS TS V4.R2 increased the number to 255
- Allows the access to different subpools simultaneously (different files) from different TCBs
- Lock mechanism is used to protect integrity
- Multiple pools for DB2/MQ CICS regions
- In case of an FOR single pool is recommended as VSAM Threadsafe is not available (FCQRONLY=YES)
- IPIC supports threadsafe Function Shipping (4.2)
 - FCQRONLY=NO)





• LSR VSAM Threadsafe files

- Allows the execution of File Control commands on Open TCBs
- As the LSR pools can be accessed by multiple TCBs simultaneously, a lock structure was developed to ease the single thread access to LSR pools
 - Separate executing tasks on different TCBs can access different LSR Pools simultaneously
 - Therefore, consideration should be given to the use of multiple pools with VSAM threadsafe to improve throughput
 - Tasks must be defined as threadsafe to use VSAM threadsafe





- LSR VSAM Threadsafe file distribution
 - Prior to CICS TS V4.R2 only 8 LSR pools are available
 - Distribution of files to obtain overlap requires planning due to the limited number of pools
 - Have one general pool (e.g., # 1) for non-threadsafe applications and one for data tables (if any) and use the remainder to distribute the files
 - If limited number of threadsafe applications, distribute the remaining pools among the threadsafe applications
 - Determine application activity and assign the remaining pools by application
 - May require more virtual/real storage





- LSR VSAM Threadsafe file distribution
 - CICS TS V4.R2 increased the number of LSR pools available 255
 - Distribution files is much easier due to increased number of pools available
 - Assign one general pool (e.g., # 1) for non-threadsafe applications and one for data tables (if any)
 - Allocate separate pools to important highly active files
 - May require more virtual/real storage
 - For FOR regions use IPIC which has a threadsafe CSMI





Buffer Fragmentation

- Only Eleven Valid CISZ for LSR Buffers (K)
 - 0.5 1.0 2.0 4.0 8.0 12.0
 - 16.0 **20.0 24.0 28.0 32.0**
- Therefore, a 2.5K Byte CISZ Would Use a 4K LSR Buffer
- If a 4K Buffer Was Not Available, Then the Next Largest Available Buffer Is Used
- Some Fragmentation May Be Desired for Certain CISZ (e.g., non VSAM/E – 18.0K)





Buffer Fragmentation

- Avoid Unnecessary Fragmentation (e.g., a 6K CISZ Using a 12K Buffer)
- Certain Default Index CISZ Should Be Forced to an LSR CISZ (e.g., 1536 to 2048 or 2560 to 4096)
- Virtual Fragmentation Results in Real Storage Fragmentation





LSR Buffer vs. File CISZ Reconciliation

- Best Alternative to Reducing Fragmentation
- Determine File CI Sizes Required and Assign LSR Pool Buffers to Match
 - Number and Size of Buffers
 - Number of Strings (Overall)
- Set CISZ Standards (If possible) for LSR Pool Files
- Complex Task, If Done Manually





LSR buffer vs. file CISZ reconciliation

- Some installations simply define a certain number of buffers for every possible buffer Size (11 buffer sizes possible in an LSR pool)
- Alternate example:
 - Suppose you don't have any 16K buffer users (CISZ range is 14K and 16K files)
 - You determine that you want to have twenty 16K buffers defined (320 K) just in case one day you get a 14K or 16K file
 - This allocated storage will not be used wasted storage every day of the year
 - Instead, why don't you simply define sixteen 20K buffers (320K) (or next useable size) that will be used every day





- Page boundary buffer allocation (Minor)
 - VSAM requests buffers on a page boundary and in page (4K) increments
 - Fragmentation that occurs from buffer allocation should be avoided – loss of virtual storage
 - Allocate the following buffers in the following increments:
 - 0.5K Multiple of 8 (0.5K Times 8 = 4K)
 - 1.0K Multiple of 4 (1.0K Times 4 = 4K)
 - 2.0K Multiple of 2 (2.0K Times 2 = 4K)





- Maximum key size (Minor)
 - Maximum key size is important as all VSAM control blocks are shared and must be able to accommodate the largest file key in the shared pool
 - If the maximum key size allocated to the pool Is too small, files with larger keys will not open
 - Many installations force the LSR pool key size to 255 bytes
 - Although using this maximum can waste storage, the actual amount wasted depends on the number of strings allocated times the excess key size
 - Decision is installation dependent





Number of strings allocated

- Probably only tuned when wait on strings conditions occur
 - String waits can occur If
 - Maximum number of strings in the pool is reached
 - Maximum number of strings assigned to the file is reached
- Many LSR pools strings are over-allocated
- The objective should be to have sufficient strings to handle peak periods without having to wait on strings
- Try to allocate strings so that the high used string number is around 50 to 60% of the total strings allocated to the pool
- Before increasing strings due to wait on strings conditions, make sure that you are attaining your look-aside hit ratio objectives for the pool





LSR Pool Candidates

- LSR provides the best look-aside algorithm within CICS
- Generally, files (high, intermediate and low activity) should be assigned to LSR except:
 - Share Options 4 files
 - Files that do not follow Command Level guidelines for accessing VSAM
 - Start Browse, Read NextRead for Update (Non-RLS)
 - High CA split activity files (tune independently)
- LSR Is the gate to new file features within CICS





LSR Recommendations

- LSR Is preferred over NSR buffering
 - Superior look-aside hit ratio
- Tuning LSR involves:
 - Ensuring proper number of buffers defined
 - Achieve installation look-aside hit ratio goals
 - Eliminating fragmentation
 - Static definition of the pool(s)
- Continuous review especially when major application changes occur
 - VSAM tuning



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Closing

- Use LSR over NSR
- Tune to eliminate I/O Look-Aside Hits
- Monitor File Statistics periodically to ensure that Look-Aside Hit Ratio objectives are being met
- When tuning LSR remember Robin Hood!

